

Day 1

6/26/11

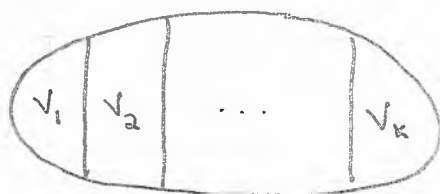
# Colorings and Homomorphisms for Graphs and Finite Structures

Jaroslav Nešetřil

Combinatorics Surrounding CSP  
Paradoxes of coloring

$G = (V, E)$   $\chi(G) = \min$  # of colors which suffice to color  $G$  s.t. adj. vertices receive diff. colors.

$G$



$v_i$  no edges "easy"  
min # of partitions into "easy" parts.

edges are "between"

$$\chi(G) \leq \max_{x \in V} d_G(x) + 1$$

$$\leq \max \text{ average degree} + 1$$

" degeneracy

$$\chi(G) \geq \omega(G) = \max \{k \mid K_k \text{ contained in } G\}.$$

Johnson  
↳ IF  $G$  contains no triangle  $\Delta$ , then  $\chi(G) \leq \frac{|\Delta(G)|}{\log_3 |\Delta(G)|}$

$$\chi(G) \leq \frac{3\Delta(G)}{4} \leq c \cdot \Delta(G)$$

## Homomorphism

$$G = (V, E) \quad G' = (V', E') \quad (\text{undirected})$$

$$f: G \rightarrow G'$$

$$f: V \rightarrow V' \quad xy \in E \Rightarrow f(x)f(y) \in E'$$

$$G = Q_d \text{ (d-dimensional cube)} \xrightarrow{f} \dots \circ \circ \circ \circ \circ \dots$$

J. Kahn: Expected range  $\leq c_0 = 5$

D. Galvin

$$G, H \quad \text{Hom}(G, H) = \{ f \mid f : G \rightarrow H \}$$

$$\text{hom}(G, H) = |\text{Hom}(G, H)|$$

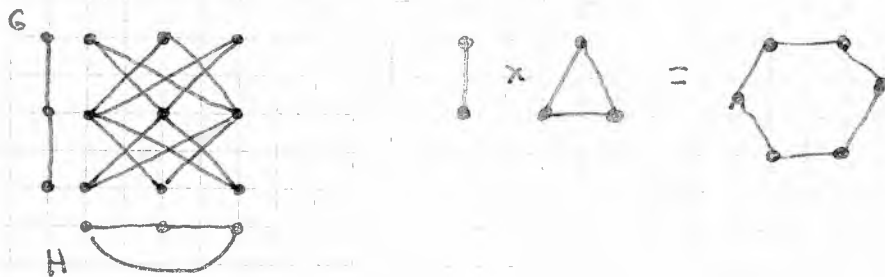
Thm  $(L^2) \quad H_1 \cong H_2 \Leftrightarrow \forall G \quad \text{hom}(G, H_1) = \text{hom}(G, H_2)$

$$t(G, H) = \frac{\text{hom}(G, H)}{|H|^{|G|}} \quad (|H| = |V(H)|)$$

← all maps

so this is probability that a map is a homomorphism

$G \times H =$  <sup>direct</sup> categorical product of  $G$  and  $H$



Thm  $\text{SP} \{ K_k \mid k=2,3,\dots \} = \text{all graphs}$

$$G \leq_{\text{ind}} \bigcup_{i \in I} K_{n_i}$$

$\leq_{\text{ind}}$  = induced subgraph.

Defn  $\text{dim}(G) = \min |I| = \min \{ d \mid G \leq_{\text{ind}} K_{|G|}^d \}$

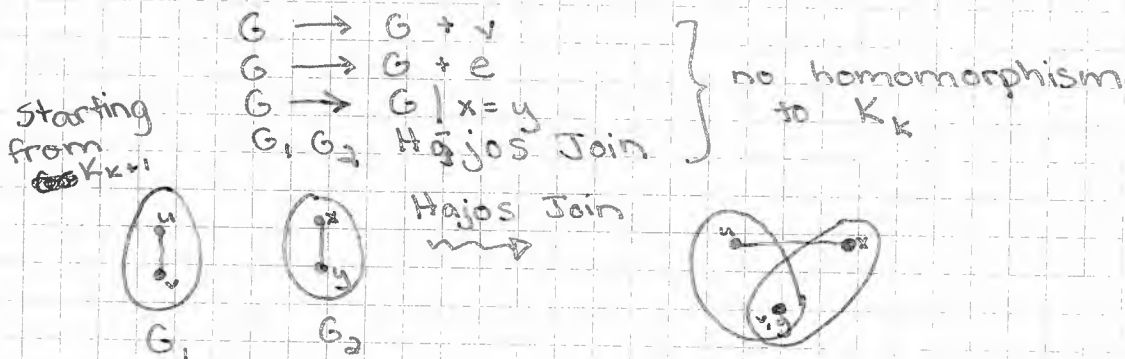
Thm  $\forall k \exists A_k \quad \chi(G) \leq k \Leftrightarrow G \leq_{\text{ind}} A_k^k$

minimal examples



⋮

Remark | Any  $G \rightarrow K_k$  can be generated by adding a vertex or edge or identifying 2 vertices



We don't know what  $\chi(G)$  large means

Conjectures

① Product conjecture (Hedetniemi)  $\chi(G \times H) = \min\{\chi(G), \chi(H)\}$

↑  
know  $\leq$

$f(k) = \min\{\chi(G \times H) \mid \chi(G), \chi(H) \geq k\}$   
 conjecture says  $f(k) = k$ . (true for  $k=2, 3, 4$ )

?  $f(k) \rightarrow \infty$  Poljak, Rodl

X. Zhu  $\chi_f(G \times H) = \min\{\chi_f(G), \chi_f(H)\}$

$\chi_f$  = fractional chromatic number

$\chi_f(G) = \min\left\{\frac{n}{k} \mid G \rightarrow K\left(\frac{n}{k}\right)\right\}$

$K\left(\frac{n}{k}\right) = (V, E)$   $V =$  all  $k$ -elt. subsets of  $\{1, \dots, n\}$   
 $E =$  pairs of disjoint subsets

↑  
Kneser Graphs

$\chi(K\left(\frac{n}{k}\right)) = n - 2k + 2$  (Lovasz?)

(Erdős, 1958)  
Thm |  $\forall k \forall g \exists G$  s.t.  $\chi(G) \geq k$  and  $G \not\cong \Delta, \square, \dots, C_g$   
 (girth  $G > g$ )

How large  $G$ ?  
 Construction of  $G$ ? means  $\chi$  is not "local"

All  $|G|$  not primitive recursive in  $k, g$

② If  $\chi(G) > f(k, g)$  large then  $\exists H \subseteq G$  s.t.  $\chi(H) \geq g$  and

$\text{girth}(H) \geq g$ . (Erdős, Hajnal)

$g=3$  true Rödl 1977