

# Omitting types in infinitary $[0, 1]$ -valued logic

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1. Background - the first-order case
2. Infinitary logic for metric structures and omitting types
3. Applications

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## Example

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 $\langle \overline{\mathbb{Q}}, +, -, \cdot, 0, 1 \rangle \not\cong \langle \overline{\mathbb{Q}(\pi)}, +, -, \cdot, 0, 1 \rangle$

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 $\langle \overline{\mathbb{Q}}, +, -, \cdot, 0, 1 \rangle \not\cong \langle \overline{\mathbb{Q}(\pi)}, +, -, \cdot, 0, 1 \rangle$

No formula, even with parameters from  $\overline{\mathbb{Q}}$ , can distinguish these structures.

## Definition

Let  $\Sigma(x_1, \dots, x_n)$  be a set of  $L$ -formulas. If  $\mathcal{M}$  is an  $L$ -structure such that there are  $a_1, \dots, a_n \in \mathcal{M}$  such that  $\mathcal{M} \models \phi(a_1, \dots, a_n)$  for all  $\phi(x_1, \dots, x_n) \in \Sigma$ , we say  $\mathcal{M}$  **realizes**  $\Sigma$ . Otherwise  $\mathcal{M}$  **omits**  $\Sigma$ .

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If  $T$  is an  $L$ -theory and there is  $\mathcal{M} \models T$  such that  $\mathcal{M}$  realizes  $\Sigma$ , then we say  $\Sigma$  is a **type of**  $T$ .

## Example

Let  $\Sigma(x) = \{a_0 + a_1x + \cdots + a_nx^n \neq 0 : n \in \omega, a_i \in \mathbb{Z}\}$ .

Then  $\pi$  realizes  $\Sigma$  in  $\overline{\mathbb{Q}(\pi)}$ , but  $\overline{\mathbb{Q}}$  omits  $\Sigma$ .



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Types increase expressive power.

## Example

Let  $T$  be the theory of abelian groups, and

$$\Sigma(x) = \{x \neq 0, x + x \neq 0, x + x + x \neq 0, \dots\}.$$

Then  $\mathcal{M} \models T$  omits  $\Sigma$  if and only if  $\mathcal{M}$  is a torsion group.

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## Example

Let  $T = PA$ , and

$$\Sigma(x) = \{x \neq 0, x \neq S0, x \neq SS0, \dots\}.$$

Then  $\mathcal{M} \models T$  omits  $\Sigma$  if and only if  $\mathcal{M}$  is standard.

## Definition

A type  $\Sigma(\bar{x})$  is **principal** over  $T$  if there is  $\phi(\bar{x})$  such that  $T \cup \{\phi(\bar{x})\} \models \Sigma(\bar{x})$ .

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## Theorem (Henkin-Orey 1957)

*For each  $n$ , let  $\Sigma_n$  be a type of  $T$  which is non-principal. Then there is a (countable)  $\mathcal{M} \models T$  which omits every  $\Sigma_n$ .*

## Corollary

*Prime models realize only principal types.*

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*Suppose  $\mathcal{M}$  is an ordered field, and that for every  $\phi$  if  $\mathcal{M} \models \exists x \phi(x)$  then there is a finite  $a \in \mathcal{M}$  such that  $\mathcal{M} \models \phi(a)$ . Then  $\mathcal{M}$  is elementarily equivalent to an Archimedean field.*

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## Corollary

*If  $\mathcal{M} \models PA$  (or  $\mathcal{M} \models ZF$ ) is countable then there exists a countable end-extension  $\mathcal{N}$  of  $\mathcal{M}$  such that  $\mathcal{M} \preceq \mathcal{N}$ .*



## Theorem (Keisler 1973)

*Let  $L$  be a countable fragment of  $\mathcal{L}_{\omega_1, \omega}$ , and let  $T$  be an  $L$ -theory. For each  $n$ , let  $\Sigma_n$  be a type of  $T$  which is non-principal. Then there is a (countable)  $\mathcal{M} \models T$  which omits every  $\Sigma_n$ .*

# $[0, 1]$ -valued Model Theory

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**Structures:** Metric spaces of diameter  $\leq 1$ . Interpretations of symbols respect the moduli of continuity from the signature.

## $[0, 1]$ -valued $\mathcal{L}_{\omega_1, \omega}$

Fix a signature  $S$ . The formulas of  $\mathcal{L}_{\omega_1, \omega}(S)$  are:

**Atomic Formulas:**  $d(x, y)$ ,  $R(x_1, \dots, x_n)$ , constants for each  $r \in \mathbb{Q} \cap (0, 1)$ .

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- ▶  $\sup_x \phi$
- ▶  $\sup_{n \in \omega} \phi_n$

We can recover other connectives as limits. Starting with:

$$\frac{1}{2}x = \lim_{n \rightarrow \infty} \bigvee_{i=1}^n \left( \frac{i}{n} \wedge \neg(x \rightarrow \frac{i}{n}) \right).$$

By Stone-Weierstrass, combinations of our connectives uniformly approximate any continuous  $F : [0, 1]^n \rightarrow [0, 1]$ .

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### Example

In the signature of Banach algebras, let  $\phi(x)$  be the formula  $1 - \|x \cdot x\|$ . Then  $\mathcal{M} \models \phi(a)$  if and only if  $a^2 = 0$  in  $\mathcal{M}$ .

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**Note:**  $\mathcal{M} \models \phi \rightarrow \psi$  if and only if  $\phi^{\mathcal{M}} \leq \psi^{\mathcal{M}}$ .

Examples of classes of Banach spaces axiomatizable in  $\mathcal{L}_{\omega_1, \omega}$ :

- ▶ Classes axiomatizable in finitary continuous logic ( $L^p(\mu)$  spaces,  $C(K)$  spaces, some classes of Nakano spaces, ...),

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- ▶ uniformly convex spaces,
- ▶ spaces which are not super-reflexive,
- ▶ spaces which are not hereditarily indecomposable,
- ▶ spaces which are unstable (in the sense of Krivine-Maurey).

A **fragment** of  $\mathcal{L}_{\omega_1, \omega}(S)$  is a set of formulas  $L$  such that:

- ▶  $L$  contains every atomic formula
- ▶  $L$  is closed under  $\rightarrow, \neg, \wedge, \vee, \sup_x$
- ▶  $L$  is closed under substituting terms for free variables
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From now on,  $L$  is a fixed countable fragment of  $\mathcal{L}_{\omega_1, \omega}(S)$ ,  $S$  has no function symbols. For  $C$  a set of new constant symbols,  $L_C$  is the least fragment of  $\mathcal{L}_{\omega_1, \omega}(S \cup C)$  containing  $L$ .

## Definition

Let  $T$  be an  $L$ -theory. A type  $\Sigma(\bar{x})$  is **principal** over  $T$  if there is a formula  $\phi(\bar{x})$  consistent with  $T$  such that for some  $r \in \mathbb{Q} \cap (0, 1)$ ,  $T \cup \{\phi(\bar{x}) \geq r\} \models \Sigma(\bar{x})$ .

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## Theorem

Let  $T$  be an  $L$ -theory, and let  $\{\Sigma_n(\bar{x}_n) : n \in \omega\}$  be a collection of non-principal types of  $T$ . Then there is a (countable)  $\mathcal{M} \models T$  omitting each  $\Sigma_n$ .

# Topology

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Note  $\mathcal{M}, \mathcal{N}$  are indistinguishable if and only if  $\mathcal{M} \equiv_L \mathcal{N}$ , so the space is not  $T_0$ .

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Say  $\mathcal{M} \notin \text{Mod}(T)$ , so some  $\phi$  is such that  $\phi^M < 1$ .

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Note  $\phi^{-1}([a, b]) = \{\mathcal{N} : a \leq \phi^{\mathcal{N}} \leq b\} = Mod(a \leq \phi \wedge \phi \leq b)$  is closed, so  $\phi$  is continuous. □

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If  $\Sigma(\bar{x})$  is principal, there is  $\phi, r$  such that  $T \cup \{\phi(\bar{x}) \geq r\} \models \Sigma(\bar{x})$ .

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Then for any  $r' \in (r, 1)$  we have

$$\text{Mod}_{L_{\bar{x}}}(T) \cap \text{Mod}_{L_{\bar{x}}}(\phi(\bar{x}) > r')$$

is a non-empty open subset of  $\text{Mod}_{L_{\bar{x}}}(T \cup \Sigma(\bar{x}))$ .



Proof (con't).

Conversely, there is  $\phi(\bar{x})$  such that

$$Mod_{L_{\bar{x}}}(T) \cap Mod_{L_{\bar{x}}}(\phi(\bar{x}) > 0)$$

is a non-empty open subset of  $Mod_{L_{\bar{x}}}(T \cup \Sigma(x))$ .

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Conversely, there is  $\phi(\bar{x})$  such that

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So there is  $r \in \mathbb{Q} \cap (0, 1)$  such that  $T \cup \{\phi(\bar{x}) \geq r\}$  is satisfiable.

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So there is  $r \in \mathbb{Q} \cap (0, 1)$  such that  $T \cup \{\phi(\bar{x}) \geq r\}$  is satisfiable.

Let  $\psi(\bar{x})$  be the formula  $\phi(\bar{x}) \geq r$ . Then  $\psi$ ,  $1 - s$  witness  
principality for any  $s \in \mathbb{Q} \cap (0, r)$ . □

We add countably many new constants  $C = \{c_n : n \in \omega\}$  to  $S$ .  
We work in  $\mathcal{W} \subseteq \text{Str}(L_C)$  where satisfaction of  $\sup_x \phi$  is witnessed  
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### Lemma

*For any  $\mathbf{i} = i_0, \dots, i_{n-1}$ , the natural map  $R_{\mathbf{i}} : \mathcal{W} \cap \text{Mod}(T) \rightarrow \text{Mod}_{L_{\mathbf{q}}}(T)$  is open, continuous, and surjective.*

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### Lemma

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surjective.*

### Proposition

*Let  $\Sigma(x_0, \dots, x_{n-1})$  be a nonprincipal type of  $T$ . Then for each  
 $\mathbf{i} \in \omega^n$ ,  $R_{\mathbf{i}}^{-1}(\text{Mod}(T \cup \Sigma(\mathfrak{c}_i)))$  is closed nowhere dense in  
 $\mathcal{W} \cap \text{Mod}(T)$ .*

Suppose there is

$$\langle \mathcal{M}, \bar{a} \rangle \in \mathcal{W} \cap \text{Mod}(T) \setminus \bigcup_{n \in \omega} \bigcup_{i \in \omega^n} R_i^{-1}(T \cup \Sigma_n(\mathcal{G}_i)).$$

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Since  $\langle \mathcal{M}, \bar{a} \rangle \in \mathcal{W}$ ,  $\bar{a} \preceq \mathcal{M}$ , and our choice ensures no subset of  $\bar{a}$  realizes any  $\Sigma_n$ , so  $\bar{a} \models T$  and omits each  $\Sigma_n$ .



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Since  $\langle \mathcal{M}, \bar{a} \rangle \in \mathcal{W}$ ,  $\bar{a} \preceq \mathcal{M}$ , and our choice ensures no subset of  $\bar{a}$  realizes any  $\Sigma_n$ , so  $\bar{a} \models T$  and omits each  $\Sigma_n$ .

It therefore suffices to show that  $\mathcal{W} \cap \text{Mod}(T)$  is **Baire**, i.e., the countable union of closed nowhere dense sets is codense.

## Definition

Let  $X$  be a completely regular space. A **complete sequence of open covers** of  $X$  is a sequence  $\langle \mathcal{U}_n : n \in \omega \rangle$  of open covers such that if  $\mathcal{F}$  is a centred family of closed sets such that for each  $n \in \omega$  there is  $F_n \in \mathcal{F}$  and  $U_n \in \mathcal{U}_n$  such that  $F_n \subseteq U_n$ , then  $\bigcap \mathcal{F} \neq \emptyset$ .

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## Definition

A completely regular space  $X$  is **Čech-complete** if it has a complete sequence of open covers.

## Fact

*If  $X$  is  $T_{3\frac{1}{2}}$  then  $X$  is Čech-complete if and only if  $X$  is a  $G_\delta$  in  $\beta X$ .*

## Fact

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## Fact

*If  $X$  is completely metrizable or locally compact Hausdorff, then  $X$  is Čech-complete.*

## Proposition

*Let  $X$  be completely regular and Čech-complete. Then:*

- ▶  *$X$  is Baire.*
- ▶ *Every closed subspace of  $X$  is Čech-complete.*

## Theorem

$\mathcal{W}$  is Čech-complete.

- Proof:** Recursively build a sequence of open covers such that:
- ▶ The sets in  $\mathcal{U}_n$  prescribe ranges for the first  $n$  sentences.



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  - ▶ i.e., if  $U \in \mathcal{U}_n$  then for each  $m \leq n$  there is  $I_m \subseteq [0, 1]$  such that for all  $\mathcal{M} \in U$ ,  $\sigma_m^{\mathcal{M}} \in I_m$ .

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- ▶ When a set prescribes a range for  $\sup_n \phi_n$  it also picks an  $n$  and specifies that  $\phi_n$  be in the same range.

**Proof:** Recursively build a sequence of open covers such that:

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Let  $\mathcal{F}$  be a centred family of closed sets. For each  $n$ , pick  $F_n \in \mathcal{F}$ ,  $U_n \in \mathcal{U}$  such that  $F_n \subseteq U_n$ . For  $F \in \mathcal{F}$ , let  $T_F$  be a theory such that  $F = \text{Mod}(T_F)$ .

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$$\bigcap \mathcal{F} = \bigcap_{n \in \omega} F_n.$$

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So  $F \cap F_n = \emptyset$ , contradiction. □

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If  $\{n \in \omega : \sigma^{M_n} > a\} \in \mathcal{D}$ , by induction it suffices to find  $m^* \in \omega$  such that  $\{n \in \omega : \theta_{m^*}^{M_n} > a\} \in \mathcal{D}$ .



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So cofinitely many  $\mathcal{M}_n \models \theta_{m^*} > a$ , so  $\{n \in \omega : \theta_{m^*}^{\mathcal{M}_n} > a\} \in \mathcal{D}$ .

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Similar proof for  $\sigma = \sup_x \phi$ . □

It follows that  $\mathcal{M} \in \bigcap_{n \in \omega} F_n = \bigcap \mathcal{F} \neq \emptyset$ . So  $\mathcal{W}$  is Čech-complete.



## Theorem

Let  $T$  be an  $L$ -theory, and let  $\{\Sigma_n(\bar{x}) : n \in \omega\}$  be a collection of non-principal types of  $T$ . Then there is  $\mathcal{M} \models T$  omitting each  $\Sigma_n$ .

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## Example

If  $(c_i)_{i \in \omega}$  are constant symbols and  $T$  is a theory which implies that  $(c_i)$  is a Cauchy sequence, then the type  $\Sigma(x)$  expressing that  $x = \lim_{i \rightarrow \infty} c_i$  can be omitted in a metric structure, but not in a complete metric structure.

## Definition

For  $\Sigma(x_1, \dots, x_n)$  a type, and  $\delta \in \mathbb{Q} \cap (0, 1)$ ,

$$\Sigma^\delta = \left\{ \sup_{y_1} \dots \sup_{y_n} \left( \bigwedge_{k \leq n} d(x_k, y_k) \leq \delta \wedge \sigma(y_1, \dots, y_n) \right) : \sigma \in \Sigma \right\}.$$



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## Definition

A type  $\Sigma(\bar{x})$  of  $T$  is **metrically principal** over  $T$  if for every  $\delta > 0$  the type  $\Sigma^\delta(\bar{x})$  is principal over  $T$ .

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*Let  $T$  be an  $L$ -theory, and let  $\{\Sigma_n(\bar{x}) : n \in \omega\}$  be a collection of types of  $T$  which are not metrically principal. Then there is  $\mathcal{M} \models T$  such that  $\overline{\mathcal{M}}$  omits each  $\Sigma_n$ .*

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## Definition

A fragment  $L$  of  $\mathcal{L}_{\omega_1, \omega}$  is **continuous** if every formula  $\phi(x_1, \dots, x_n) \in L$  defines a *continuous* function  $\phi : \mathcal{M}^n \rightarrow [0, 1]$  for every  $L$ -structure  $\mathcal{M}$ .

## Theorem

*Let  $T$  be an  $L$ -theory, where  $L$  is a countable continuous fragment, and let  $\{\Sigma_n(\bar{x}) : n \in \omega\}$  be a collection of types of  $T$  which are not metrically principal. Then there is a complete  $\mathcal{M} \models T$  omitting each  $\Sigma_n$ .*

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- ▶ If  $L$  is first-order and we restrict to discrete structures, this is the classical Omitting Types Theorem.

# Application - Separable Quotients

## Question

*Let  $X$  be a Banach space. Is there a subspace  $Y \subseteq X$  such that  $X/Y$  is separable?*

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Equivalently,

## Question

*Let  $X$  be a non-separable Banach space. Is there a separable Banach space  $Y$  and a surjective bounded linear operator  $T : X \rightarrow Y$ ?*

## Theorem

*Let  $X, Y$  be Banach spaces with  $\text{density}(X) > \text{density}(Y)$ , and let  $T : X \rightarrow Y$  be a surjective bounded linear map.*

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- ▶  $(X, Y, T) \equiv_L (X', Y', T')$ .

## Proof (Outline).

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Using Downward Löwenheim-Skolem, let  
 $\langle X_0, Y_0, T_0, L_0, c, \leq_0, f_0 \rangle = \mathcal{M}_0 \preceq \mathcal{M}$  be countable.

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Repeat  $\omega_1$  times, and take the union of the elementary chain to get  $\mathcal{M}_{\omega_1} = \langle X_{\omega_1}, Y_{\omega_1}, T_{\omega_1}, \dots \rangle$ .

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Take  $\langle X', Y', T' \rangle$  to be the completion of  $\langle X_{\omega_1}, Y_{\omega_1}, T_{\omega_1} \rangle$

## Example

Suppose  $T : X \rightarrow Y$  is bounded, linear, surjective, where  $\text{density}(X) > \text{density}(Y)$ . Then there are  $X', Y'$  and  $T' : X' \rightarrow Y'$ , with  $\text{density}(X') = \aleph_1$  and  $Y'$  separable, such that:

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- ▶  $Y$  and  $Y'$  have isometric ultrapowers,
- ▶ if  $X$  is uniformly convex then so is  $X'$ ,
- ▶ if  $X$  is not hereditarily indecomposable then neither is  $X'$ ,

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Thank you!